## Solutions 9-10

9.1 Show that removing a small disc from a projective plane we get a Möbius band.

Solution. One can use the following "cut and paste" argument:

9.2 Removing a line from $\mathbb{R}^{2}$ or from $S^{2}$ one gets a space with two connected components. Show that $R P^{2} \backslash \mathbb{R} P^{1}$ is a connected space.

Solution. Recall that $\mathbb{R} P^{2}$ may be obtained obtained from $S^{2}$ by identification of antipodal points. Let as remove the line before the identification: we will get two open hemispheres. After the identification we get an open disc.
9.3. Let the Möbius band $\mathbf{M}$ is obtained by gluing along the vertical sides of the square with vertices $( \pm 1, \pm 1)$. Let $m$ be a midline of the Möbius band $\mathbf{M}$ (obtained from the segment of the line line $y=0$ ).
(a) what is $\mathbf{M} \backslash m$ ?
(b) Let $l$ be the closed line obtained from $y=1 / 2$ and $y=-1 / 2$. What is $\mathbf{M} \backslash l$ ?

Solution. Glue the Möbius band from a piece of paper, take the scissors and cut along $m$ or $l$ respectively.
9.4. Let $\mathbf{C}$ be the conic $x^{2}+y^{2}=z^{2}$. What kind of space is $\mathbb{R} P^{2} \backslash \mathbf{C}$ ?

Solution. Looking at a hemisphere it is easy to see that $\mathbb{R} P^{2} \backslash \mathbf{C}$ has two connected component, one is the disc ("inside" of the conic), the other is a projective plane without disc, which is a Möbius band as shown in Question 9.1.
9.5. Given a point $P$ inside a circle and a chord $A B$ through the point $P$, let $M_{A B}$ denote the intersection point of the two lines tangent to the circle at $A$ and $B$. Show that $M_{A B}$ runs over some line as $A$ runs over the circle.

Solution. When $A$ runs over a circle, $M_{A B}$ runs over all points polar to $A$, so we get a line polar to $A$.
9.6. Let $\mathbf{C}$ be the conic $x^{2}+y^{2}=z^{2}$. A triangle in $\mathbb{R} P^{2}$ is self-polar (with respect to $\mathbf{C}$ ) if its sides are polar to its vertices (not necessarily the opposite ones).
(a) Construct a self-polar triangle with two vertices on C.
(b) Does there exist a self-polar triangle with exactly one vertex on $\mathbf{C}$ ?
(c) Show that there exists a self-polar triangle having no vertex on $\mathbf{C}$. Hint: it may have some vertices at infinity.

Solution. (a) A point $A \in \mathbf{C}$ is polar to a line $l_{A}$ tangent to $\mathbf{C}$ at $A$. So, if we take a point $C$ outside the conic, and consider the tangent lines $L_{A}$ and $L_{B}$ through $C$ (tangent to $\mathbf{C}$ at $A$ and $B$ respectively), then $c=A B$ is polar to $C, A$ is polar to $a$ and $B$ is polar to $b$.
(b) Suppose that such a configuration $A B C$ does exist. Let $A$ be the vertex on $\mathbf{C}$ and $a$ the tangent at $A$. The side $a$ of the triangle should have two vertices on it, one is $A$ let $B$ be the other one. Next, let $b$ be the side of the triangle polar to $B$ : it should pass through $A$ and another point where a line through $B$ is tangent to $\mathbf{C}$. The last vertex $C$ of the triangle $A B C$ should lie on the intersection of $b$ and the line $c$ polar to $C$. By assumption, $C \notin \mathbf{C}$. If $C$ lies inside the conic then the polar line $c$ lies completely outside the conic, so $C \notin c$. If $C$ lies outside the conic, then $c$ intersects $b$ inside the conic, so $C \neq b \cap c$.
(c) One can take a triangle with $A$ at the centre of the circle, and $B$ and $C$ at infinity, so that the lines $b$ and $c$ form a right angle and line $a$ is the line at infinity.

(a)

(c)



Remark. If you want to see such a configuration without vertices at infinity, look at Figure (d). It rather easy to see that this configuration does exist (from continuity reasons), and more difficult (but still possible) to find an explicit construction.
9.7. (a) Formulate the theorem dual to Desargues' theorem.
(b) Draw an example. (Hint: send the line $s$ to the line at infinity).
(c) Can you prove this theorem?

Solution. (a) We need to substitute: points by lines, lines by points, a line through two points by a point of intersection of two lines, intersection of two lines by a line through two points, concurrency of lines by collinearity of points (and backwards).
The assumption was:
Suppose that the lines joining the corresponding vertices of triangles $A_{1} A_{2} A_{3}$ and $B_{1} B_{2} B_{3}$ intersect at one point $S$.

Denote by $a_{i}$ and $b_{i}$ the sides of the triangles opposite to $A_{i}$ and $B_{i}$ respectively. Then the assumption says:

Suppose that the lines $A_{1} B_{1}, A_{2} B_{2}, A_{3} B_{3}$ intersect at one point $S$.
We rewrite:
Suppose that the points $a_{1} \cap b_{1}, a_{2} \cap b_{2}, a_{3} \cap b_{4}$ lie on one line $s$.
The conclusion was
Then the points $P_{1}=A_{2} A_{3} \cap B_{2} B_{3}, P_{2}=A_{1} A_{3} \cap B_{1} B_{3}, P_{3}=A_{1} A_{2} \cap B_{1} B_{2}$ are collinear.
Instead of points $P_{i}$ we will get lines $p_{i}$ :
$p_{1}$ is a line through $a_{2} \cap a_{3}$ and $b_{2} \cap b_{3}$, in other words, $p_{1}=A_{1} B_{1}$, and similarly, $p_{2}=A_{2} B_{2}$, $p_{3}=A_{3} B_{3}$.
So, we get the following conclusion:
Then the lines $p_{1}=A_{1} B_{1}, p_{2}=A_{2} B_{2}, p_{3}=A_{3} B_{3}$. are concurrent.
So, the dual to the Desargues' Theorem is just the converse statement:
Suppose that the points $a_{1} \cap b_{1}, a_{2} \cap b_{2}, a_{3} \cap b_{3}$ lie on one line $s$.
Then the lines $p_{1}=A_{1} B_{1}, p_{2}=A_{2} B_{2}, p_{3}=A_{3} B_{3}$. are concurrent.
(b)... and we actually don't need to draw anything new!
(c) If we send the line $s$ to infinity, then we get two triangles with corresponding sides parallel: $a_{i}$ is parallel to $b_{i}$. Let $S=A_{1} B_{1} \cap A_{2} B_{2}$. Then it is easy to see that there exists a homothety with centre $S$ which takes points $A_{1}$ to $B_{1}, A_{2}$ to $B_{2}$. It also takes the line $A_{1} A_{3}$ to the line parallel to it and passing through $A_{1}$, i.e. to $B_{1} B_{3}$. Similarly, it takes the line $A_{2} A_{3}$ to $B_{2} B_{3}$. This implies that $A_{3}=A_{1} A_{3} \cap A_{2} A_{3}$ is mapped $B_{3}=B_{1} B_{3} \cap B_{2} B_{3}$. Hence $B_{3}$ lies on the line $S A_{3}$, i.e. $S \in A_{3} B_{3}$, and te lines $A_{1} B_{1}, A_{2} B_{2}, A_{3} B_{3}$ are concurrent at $S$.
10.1. (a) How many non-intersecting lines can you draw in the Klein model of hyperbolic plane?
(b) The same question, but no other line should intersect more than two lines of your family.

Solution. (a) Choosing the segments of parallel (in Euclidean sense) lines we may get as as many mutually non-intersecting hyperbolic likes as we want (even uncountably many).
(b) To get countably many lines in the family we proceed as follow:

1) Take two disjoint lines and color the two disjoint circular segments bounded by these lines;
2) In the uncolored domain insert a line between each two lines and color the respective circular segment;
3) return to Step 2).

Here is the diagram for the first three iterations:


Now, suppose that a line $l$ intersects two lines of the family. Then both ends of $l$ should lies in the corresponding colored circular segments. It is clear that such a line will not intersect any other line in our family.

It is left to prove that it is impossible to find an uncountably big family of lines. First, suppose that in our family there are three lines $l_{1}, l_{2} a n d l_{3}$ such that the line $l_{3}$ is contained in the union of the disjoint circular segments bounded by $l_{1}$ and $l_{2}$. Then it is easy to construct a line $l$ intersecting all three lines $l_{1}, l_{2}$ and $l_{3}$. This means that in our family all lines come with the corresponding circular segment (and all these circular segments are mutually disjoint). As every circular segment contains a boundary point which makes a rational angle with horizontal line, we see that the number of the segments in our family is countable (here we use countability of rational numbers).
10.2. Show that any two lines on hyperbolic plane either intersect inside the hyperbolic plane, or intersect on its boundary, or have a unique common perpendicular.

Solution. Let $A B$ and $C D$ be two lines having no intersection neither inside nor on the boundary of the hyperbolic plane. Let $l_{A}, l_{B}, l_{C}$, and $l_{D}$ be the lines tangent to the disc. A line orthogonal to $A B$ in the Klein model should pass through the intersection point $X=l_{A} \cap l_{B}$. Similarly, the line orthogonal to $C D$ should pass through $Y=l_{C} \cap l_{D}$. There exists a unique line on (Euclidean plane) through $X$ and $Y$, it gives the common perpendicular. See on the diagram:

10.3. (Right-angled polygons on hyperbolic plane)
(a) Show that a hyperbolic triangle can not have more than 1 right angle.
(b) Show that there are no hyperbolic rectangles (i.e. quadrilaterals with 4 right angles).
(c) In the Klein model, construct a hyperbolic pentagon with 5 right angles.

Solution. (a) Let $C$ be a right-angled vertex of the triangle, without loss of generality (i.e. by transitivity of isometry group on the points of the model) we may assume that $C$ is the centre of the disc. The two sides $a$ and $b$ through $C$ will be two orthogonal (in Euclidean sense) diameters. Now, a hyperbolic line orthogonal to a diameter (in hyperbolic sense) is also orthogonal to it in Euclidean sense. So if $c$ is orthogonal to $a$ then it will never meet $b$ inside the model, and we can not get a triangle.
(b) Similarly to above, let $a$ and $b$ be two orthogonal sides meeting at the centre of the disc (say, $a$ horizontal, $b$ vertical). Then if $a b c d$ is a rectangle, $c$ should be horizontal and $d$ vertical. But it is easy to see that in this case $c$ and $d$ does not form a right angle ( $d$ can not pass through the intersection point $X$ of the tangents to the ends of $c$ as $X$ belongs to (continuation of) $b$ and $d$ is parallel to $b$ ).
(c) Using the same four lines as above ( $a$ horizontal diameter, $b$ vertical diameter, $c$ horizontal, $d$ vertical) we can place $c$ and $d$ far enough from the centre so that $c \cap d=\emptyset$ in hyperbolic plane. Then by Question 10.2 there exists a common perpendicular to $c$ and $d$ (and we know how to construct it):


