



Introduction to reconstructing the topological monoid of endomorphisms of the rationals.

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Presenting joint work with...

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Outline

Topological Monoids.

Automatic homeomorphicity

Clones



Reconstruction of Topology

Whether we can reconstruct the canonical topology of an **endomorphism monoid** $\text{End}(\mathbb{A})$ from its underlying abstract **monoid** structure?

Automatic continuity

In which situations are homomorphisms or isomorphisms between **transformation monoids** automatically continuous?



Reconstruction of Topology

Whether we can reconstruct the canonical topology of a **polymorphism clone** $\text{Pol}(\mathbb{A})$ from its underlying abstract **clone** structure?

Automatic continuity

In which situations are homomorphisms or isomorphisms between **function clones** automatically continuous?



Transformation monoids

- For a set A , we denote by $O_A^{(1)} := A^A$ the set of all unary functions on A and by

$$\text{Tr}(A)$$

the full transformation monoid on A .

- The submonoids

$$M \leq \text{Tr}(A)$$

are transformation monoids on A .



If we equip A with the discrete topology, then $\text{Tr}(A)$ is a product space of A equipped with the **Tychonoff topology**.

Pointwise convergence topology

Let I be an index set. For every finite $J \subseteq I$ and $u : J \rightarrow A$:

$$U(J, u) := \{f : I \rightarrow A \mid f \upharpoonright_J = u\}.$$

A basis for the topology of A^I can be expressed as

$$B_{\text{pwc}} = \{\emptyset\} \cup \{U(J, u) \mid J \subseteq I \text{ finite} \wedge u \in A^J\}.$$



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Special case $I = A$, $J = \{a_1^1, \dots, a_1^m\}$, and we fix m elements $a_0^j = u(a_1^j) \in A$ for $1 \leq j \leq m$.



Topology on $\text{Tr}(A)$

A non-empty basic open set is:

$$U(J, u) = \left\{ f: A \rightarrow A \mid \forall 1 \leq j \leq m: f(a_1^j) = u(a_1^j) = a_0^j \right\}.$$

- **Topological monoids** are abstract monoids which carry a topology under which the composition is continuous.
- A **transformation monoid** $M \leq \text{Tr}(A)$ is considered as a **topological subspace** of $\text{Tr}(A)$.



Given a relational structure $\mathbb{A} = \left(A, \left(R^{\mathbb{A}} \right)_{\underline{R} \in \Sigma} \right)$, where $R^{\mathbb{A}} \subseteq A^{\text{ar}(\underline{R})}$ for each $\underline{R} \in \Sigma$.

Endomorphism monoids

A function $f \in O_A^{(1)}$ is called an **endomorphism** of \mathbb{A} if

$$f : \mathbb{A} \xrightarrow{\text{homo}} \mathbb{A}.$$

The set of all **endomorphisms** on \mathbb{A} is denoted by

$$\text{End}(\mathbb{A}).$$



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Polymorphism

A function $f \in O_A^{(k)} := A^{A^k}$ is called a **polymorphism** of \mathbb{A} if

$$f : \mathbb{A}^k \xrightarrow{\text{homo}} \mathbb{A}.$$

The set of all **polymorphisms** on \mathbb{A} is denoted by

$$\text{Pol}(\mathbb{A}) = \bigcup_{k \in \mathbb{N}_+} \text{Pol}^{(k)}(\mathbb{A}).$$



$$f \in \mathcal{O}_A^{(k)}, \text{ar}(R^{\mathbb{A}}) = m$$

$$f \circ \left(\underbrace{\begin{pmatrix} a_{11} \\ \vdots \\ a_{m1} \end{pmatrix}}_{R^{\mathbb{A}}}, \dots, \underbrace{\begin{pmatrix} a_{1k} \\ \vdots \\ a_{mk} \end{pmatrix}}_{R^{\mathbb{A}}} \right) = \begin{pmatrix} f(a_{11}) & \dots & a_{1k} \\ & \ddots & \\ f(a_{m1}) & \dots & a_{mk} \end{pmatrix}$$



Topological closure

Remark

The submonoid $M \leq \text{Tr}(A)$ is **closed** $\iff M = \text{End}(\mathbb{A})$ for some relational structure \mathbb{A} with domain A .



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Definition (M.Bodirsky, M.Pinsker, A.Pongrácz)

A **closed** monoid $M \leq \text{Tr}(A)$ has **reconstruction** : \iff for every other closed monoid $M' \leq \text{Tr}(B)$, if there exists a monoid isomorphism between M and M' , then there also exists a monoid isomorphism between M and M' which is a homeomorphism.

Definition

A **closed** monoid $M \leq \text{Tr}(A)$ has **automatic continuity** : \iff every monoid homomorphism from M into $\text{Tr}(A)$ is continuous.

Corollary (D. Lascar (1991))

Any continuous isomorphism between closed subgroups of \mathbb{S}_A is a homeomorphism.



Definition (M.Bodirsky, M.Pinsker, A.Pongrácz)

A **closed** monoid $M \leq \text{Tr}(A)$ has **automatic homeomorphicity** : \iff every monoid isomorphism from M to a closed $M' \leq \text{Tr}(B)$ is a homeomorphism.

Some monoids with automatic homeomorphicity:

$$\text{Emb}(\mathbb{N}, =), \text{Emb}(\Gamma), \text{End}(\Gamma)$$

where $\Gamma = \text{Random graph}$.



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For groups, automatic continuity implies automatic homeomorphicity

Let \mathbb{A}, \mathbb{B} be countable. If $\text{Aut}(\mathbb{A})$ has S.I.P., then

$$\xi : \text{Aut}(\mathbb{A}) \rightarrow \text{Aut}(\mathbb{B})$$

is a homeomorphism.



We want to investigate the automatic homeomorphicity of

$\text{End}(\mathbb{Q}, \leq)$ $\text{End}(\mathbb{Q}, <)$



We want to investigate the automatic homeomorphicity of

$\text{End}(\mathbb{Q}, \leq)$ $\text{End}(\mathbb{Q}, <)$ $\text{Pol}(\mathbb{Q}, <)$ & $\text{Pol}(\mathbb{Q}, \leq)$

Constants

- For $d \in \mathbb{Q}$

$$c_d \in E := \text{End}(\mathbb{Q}, \leq)$$

where $c_d(x) := d$.

- An element $c \in M \leq O_A^{(1)}$ is called a **constant** : \iff

$$\forall x, y \in A : c(x) = c(y).$$

- The set $C = \{g \in E : (\forall f \in E) gf = g\}$ is a definable subset of E .



Proposition (M.Bodirsky, M.Pinsker, A.Pongrácz)

Let \mathbb{A} be a structure such that $\text{Pol}(\mathbb{A})$ contains all constant functions, and $\xi : \text{Pol}(\mathbb{A}) \rightarrow \mathcal{D}$ be a clone isomorphism to a clone of functions \mathcal{D} . Then ξ is open.



Proposition

Let \mathbb{A} be a structure such that $\mathcal{M}_A := \text{End}(\mathbb{A})$ contains all unary constant operations, and $\xi : \mathcal{M}_A \rightarrow \mathcal{M}_B := \text{End}(\mathbb{B})$ be a monoid isomorphism. Then ξ is open.



Proposition

Let \mathbb{A} be a structure such that $\mathcal{M}_A := \text{End}(\mathbb{A})$ contains all unary constant operations, and $\xi : \mathcal{M}_A \rightarrow \mathcal{M}_B := \text{End}(\mathbb{B})$ be a monoid isomorphism. Then ξ is open.

Let $a, b \in A$, and $E_{a,b} = \{f \in \mathcal{M}_A \mid f(a) = b\} = \{f \in \mathcal{M}_A \mid f \circ c_a = c_b\}$ be a basic open set. Then, we show that $\xi(E_{a,b})$ is open

$$\begin{aligned}\xi(E_{a,b}) &= \{\xi(f) \mid f \in \mathcal{M}_A \wedge f \circ c_a = c_b\} \\ &= \{\xi(f) \mid f \in \mathcal{M}_A \wedge \xi(f \circ c_a) = \xi(c_b)\} && \text{(since } \xi \text{ is inj.)} \\ &= \{\xi(f) \mid f \in \mathcal{M}_A \wedge \xi(f) \circ \xi(c_a) = \xi(c_b)\} && \text{(since } \xi \text{ is a hom.)} \\ &= \{g \in \mathcal{M}_B \mid g \circ \xi(c_a) = \xi(c_b)\} && \text{(since } \xi \text{ is surj.)}\end{aligned}$$

Example

$$A = \{0, 1\}$$

$$\mathcal{M}_A := \{\text{id}_A, c_0, c_1\}$$

$$c_0(x) = 0$$

$$c_1(x) = 1$$

$$B = \mathbb{N}$$

$$\mathcal{M}_B := \{\text{id}_B, e_0, e_1\}$$

$$e_0(x) = \begin{cases} 0 & \text{if } x \equiv 0 \pmod{2} \\ 1 & \text{if } x \equiv 1 \pmod{2} \end{cases}$$

$$e_1(x) = \begin{cases} 2 & \text{if } x \equiv 0 \pmod{2} \\ 3 & \text{if } x \equiv 1 \pmod{2} \end{cases}$$

- $\xi : \mathcal{M}_A \rightarrow \mathcal{M}_B$ does not map constants to constants.
- $U = \{g \in \mathcal{M}_A \mid g(0) = 0\} = \{\text{id}_A, c_0\}$ and $\xi[U] = \{\text{id}_B, e_0\}$ are basic open sets in \mathcal{M}_A and \mathcal{M}_B , respectively.



Lemma

Let $S \leq \langle A^A, \circ \rangle$ and $T \leq \langle B^B, \circ \rangle$ be *transformation semigroups* and $\xi: S \rightarrow T$ be a semigroup homomorphism, whose $\text{im}(\xi) \leq T$ *acts transitively on B* (by evaluation). That is, for all $x, y \in B$ there exists some $f_{x,y} \in S$ such that $\xi(f_{x,y})(x) = y$. In these circumstances ξ maps any constant operation $c \in S$ to a constant operation on B .



Lemma

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Proof.

- If $c \in S$ is constant, $\implies c \circ f = c$ for all $f \in S$.
- For $x, y \in B$: $c \circ f_{x,y} = c$.
- For $x, y \in B$: $\xi(c) \circ \xi(f_{x,y}) = \xi(c \circ f_{x,y}) = \xi(c)$.
- Evaluating at $x \in B$: $\xi(c)(x) = \xi(c) \xi(f_{x,y})(x) = \xi(c)(y)$,
 $\implies \xi(c)$ is a constant function.





Corollary

Let $S \leq \langle A^A, \circ \rangle$, $T \leq \langle B^B, \circ \rangle$ and $\xi: S \rightarrow T$ be a semigroup homomorphism. Suppose **S contains at least one constant operation**, then the following facts are equivalent:

- 1 $\text{im}(\xi) \leq T$ acts transitively on B (by evaluation).
- 2 $\text{im}(\xi)$ contains all unary constants on B .



Lemma

Assume,

- $\mathcal{C} \leq \langle A^A, \circ \rangle$ contains all constant operations,
- $\mathcal{D} \leq \langle B^B, \circ \rangle$ acts *transitively*,
- $\xi: \mathcal{C} \rightarrow \mathcal{D}$ *semigroup isomorphism*,

then the image of any open subset of \mathcal{C} under ξ is *open* in B^B .



Proof.

Let $a, b \in A$, and $E_{a,b} = \{f \in \mathcal{C} \mid f(a) = b\} = \{f \in \mathcal{C} \mid f \circ c_a = c_b\}$ be a basic open set. Then, we show that $\xi(E_{a,b})$ is open

$$\begin{aligned}\xi(E_{a,b}) &= \{\xi(f) \mid f \in \mathcal{C} \wedge f \circ c_a = c_b\} \\ &= \{\xi(f) \mid f \in \mathcal{C} \wedge \xi(f \circ c_a) = \xi(c_b)\} && \text{(since } \xi \text{ is inj.)} \\ &= \{\xi(f) \mid f \in \mathcal{C} \wedge \xi(f) \circ \xi(c_a) = \xi(c_b)\} && \text{(since } \xi \text{ is a hom.)} \\ &= \{g \in \mathcal{D} \mid g \circ \xi(c_a) = \xi(c_b)\} && \text{(since } \xi \text{ is surj.)} \\ &= \{g \in \mathcal{D} \mid g \circ (c_p) = c_q\}. \\ &\quad \text{(for some constants } p, q \in B, \text{ according to Lemma 5)} \\ &= E_{p,q}.\end{aligned}$$





Lemma

Assume,

- $\text{Const}_A^1 \subseteq \mathcal{C} \leq \langle A^A, \circ \rangle$
- $\mathcal{D} \leq \langle B^B, \circ \rangle$ acts *transitively*
- $\xi: \mathcal{C} \rightarrow \mathcal{D}$ *semigroup isomorphism*,

then ξ is *continuous*.

Corollary

Assume,

- $\text{Const}_A^1 \subseteq \mathcal{C} \leq \langle A^A, \circ \rangle$
- $\mathcal{D} \leq \langle B^B, \circ \rangle$ acts *transitively*
- $\xi: \mathcal{C} \rightarrow \mathcal{D}$ *semigroup isomorphism*,

then ξ is a *homeomorphism*, moreover, both \mathcal{C} and \mathcal{D} , contain all constant respective unary operations.



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Clone

$F \subseteq \mathcal{O}_A := \bigcup_{k \in \mathbb{N}_+} \mathcal{O}_A^{(k)}$ is a **clone (of operations)** on A iff

- 1 $J_A \subseteq F$
- 2 F is closed w.r.t. composition.

Definition

A function $\xi : F \rightarrow F'$ is a **clone isomorphism** iff

- 1 ξ is a bijection,
- 2 ξ respects the arities, i.e. $\text{ar}(\xi(f)) = \text{ar}(f)$ for all $f \in F$.
- 3 $\forall 1 \leq j \leq n; \xi(\pi_j^{(n)}) = \pi_j^{(n)} \in F'$,
- 4 for $f \in F^{(k)}, g_1, \dots, g_k \in F^{(m)}$ we have

$$\xi(f \circ (g_1, \dots, g_k)) = \xi(f) \circ (\xi(g_1), \dots, \xi(g_k))$$



Examples

- 1 J_A Clone of all projections.
- 2 O_A Clone of all operations.
- 3 Arbitrary intersections of clones are clones again.
Let $F \subseteq O_A$. The clone generated by F is

$$\langle F \rangle_{O_A} := \bigcap \{C \text{ is clone} \mid F \subseteq C\}$$

and it is the smallest clone containing F .

- 4 $\text{Pol}(\mathbb{A})$ for some relational structure (\mathbb{A}) .



If we equip A with the discrete topology, then $O_A^{(n)}$ is a product space of A equipped with the **product topology**.

Pointwise convergence topology

Let I be an index set. For every finite $J \subseteq I$ and $u : J \rightarrow A$:

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A basis for the topology of A^I can be expressed as

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Special case $I = A^n$, $J = \left\{ (a_1^1, \dots, a_n^1), \dots, (a_1^m, \dots, a_n^m) \right\}$, and we fix m elements $a_0^j = u(a_1^j, \dots, a_n^j) \in A$ for $1 \leq j \leq m$.



Lemma (ϕ is open)

Assume,

- $\text{Const}_A^1 \subseteq \mathcal{C} \leq \mathcal{O}_A$,
- $\mathcal{D} \leq \mathcal{O}_B$,
- $\phi: \mathcal{C} \rightarrow \mathcal{D}$ *clone isomorphism*,
- $\xi := \phi \upharpoonright_{\mathcal{C}^{(1)}}$ *semigroup homomorphism*, such that *$\text{im}(\xi)$ acts transitively* on B .

Then, for all $n > 0$ $\phi[U]$ is *open* in B^{B^n} for all open $U \subseteq \mathcal{C}^{(n)}$.



Lemma (ϕ maps any n -ary constant to an n -ary constant)

Assume,

- $\text{Const}_A^1 \subseteq \mathcal{C} \leq \mathcal{O}_A$,
- $\mathcal{D} \leq \mathcal{O}_B$, clone
- $\phi: \mathcal{C} \rightarrow \mathcal{D}$ clone isomorphism,

Then, the restriction $\xi := \phi \upharpoonright_{\mathcal{C}^{(1)}}$ maps unary constants to unary constants and ϕ maps any n -ary constant to an n -ary constant

Proof

$f \in \mathcal{O}_A^{(1)}$ constant,

$$\begin{aligned} &\iff \forall x, y \in A, f(x) = f(y) \iff f \circ \pi_1^{(2)} = f \circ \pi_2^{(2)}. \text{ Hence,} \\ &\xi(f) \circ \pi_1^{(2)} = \xi(f) \circ \xi(\pi_1^{(2)}) = \xi(f \circ \pi_1^{(2)}) = \xi(f \circ \pi_2^{(2)}) = \xi(f) \circ \pi_2^{(2)}. \\ &\implies \xi(f) \text{ is constant on } B. \end{aligned}$$



Lemma (ϕ is open)

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- $\mathcal{D} \leq \mathcal{O}_B$,
- $\phi: \mathcal{C} \rightarrow \mathcal{D}$ *clone isomorphism*,
- $\xi: \mathcal{C}^{(1)} \rightarrow \mathcal{D}^{(1)}$ *is the restriction of ϕ to the unary part of the clones and monoid isomorphism.*

Then, ξ is open and ϕ is open.

From last lemma we know

$$\phi: \text{Pol}(\mathbb{Q}, \leq) \xrightarrow{\text{clone iso.}} \mathcal{D} \text{ is open}$$



We can apply the automatic homeomorphicity of $\text{End}(\mathbb{Q}, <)$ and following proposition to show that

$$\xi : \text{Pol}(\mathbb{Q}, <) \xrightarrow{\text{clone iso.}} \mathcal{C}' \text{ is open}$$



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Proposition (32 BPP)

- Let \mathcal{C} be a topological clone on A (with the product topology) such that $\mathcal{C}^{(1)}$ acts transitively on A ,
- let ξ be *an injective clone homomorphism* from \mathcal{C} to a topological clone \mathcal{C}' (on another set B).

Suppose that the restriction $\xi \upharpoonright_{\mathcal{C}^{(1)}} : \mathcal{C}^{(1)} \rightarrow \mathcal{C}'^{(1)}$ is *open*, then so is ξ .



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Proposition (32 BPP)

- Let \mathcal{C} be a topological clone on A (with the product topology) such that $\mathcal{C}^{(1)}$ acts transitively on A ,
- let ξ be a clone isomorphism from \mathcal{C} to a topological clone \mathcal{C}' (on another set B).

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Lemma (ϕ is continuous)

Assume,

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- $\mathcal{D} \leq \mathcal{O}_B$,
- $\phi: \mathcal{C} \rightarrow \mathcal{D}$ *clone isomorphism*,
- $\xi := \phi \upharpoonright_{\mathcal{C}^{(1)}}$ *semigroup isomorphism*, and suppose $\mathcal{D}^{(1)} = \text{im}(\xi)$ *acts transitively* on B .

Then, for all $n > 0$ $\phi^{-1}[U]$ is *open* in A^{A^n} for all open $U \subseteq \mathcal{D}^{(n)}$, i.e. ϕ is *continuous*.



John K. Truss talk

$$\theta : E := \text{End}(\mathbb{Q}, \leq) \longrightarrow \text{Tr}(\Omega)$$

may be viewed as semigroup action of E on Ω .

$$\Omega = \bigcup_{i \in I} \Omega_i$$

where $\Omega_i = \{a_B^i \mid B \in [\mathbb{Q}]^{n_i}\}$, $[\mathbb{Q}]^{n_i} := \{A \subseteq \mathbb{Q} \mid |A| = n_i\}$

- $\theta(g)(a_B^i) = a_{gB}^i$ if $g \in G := \text{Aut}(\mathbb{Q}, \leq)$
- $\theta(f)(a_B^i) = a_{fB}^i$ if $f \in M := \text{Emb}(\mathbb{Q}, \leq)$
 $f \in E := \text{End}(\mathbb{Q}, \leq)$ with $|fB| = B$.



θ is continuous and open

$\theta : \text{End}(\mathbb{Q}, <) \xrightarrow{\text{inj.}} M' \leq \text{Tr}(\Omega)$ is **homeomorphism**.



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Thank you :)