

Free homogeneous structures are generalised measurable

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Thanks and acknowledgements

Thank you for giving me this chance to speak at this wonderful conference.



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This work is distinct from, but closely related to, two ongoing collaborations: one with **Dugald Macpherson**, **Charlie Steinhorn**, and **Daniel Wolf**; and the other with **Charlotte Kestner**. These collaborations have certainly informed this work and I'm greatly indebted to those colleagues **and others**. Also thanks to EPSRC.

Afterwards, I'll post these slides on anscombe.sdf.org/research.html.

Apologies for using beamer, but I will give several cumbersome definitions.

Two themes and one theorem

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- built from amalgamation of free amalgamation classes
 - e.g. random graph, generic K_n -free graphs

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Main Theorem

Free homogeneous structures in finite relational languages are generalised measurable.

Coming soon to a whiteboard near you...

1 **Theme B:** Generalised measurable structures

- Finite fields
- Measurable structures
- Pseudofinite fields, the random graph
- Generalised measurable structures

2 **Theme A:** Free homogeneous structures

- Definition
- From now on...

3 **Proof** of Main Theorem

- Definition of Γ_r
- The measuring semiring T_r
- Finishing the proof

Conventions

- 1 I will try to use the colour **magenta** for ideas that deserve a picture.
- 2 'Definable':='definable with parameters'.
- 3 ' \emptyset -definable':='definable without parameters'.
- 4 $\text{Def}(\mathcal{M})$ is the set of definable sets in the structure \mathcal{M} .
- 5 $0 \in \mathbb{N}$.
- 6 Tuples will usually be written in lowercase.
- 7 $\phi(\mathcal{M})$ is the set defined in a structure \mathcal{M} by the formula $\phi(x)$.

Theorem (Chatzidakis–van den Dries–Macintyre)

Let $\phi(x; y)$ be a formula in the language of rings. Then there exist $C \in \mathbb{R}_{>0}$, $\mu_1, \dots, \mu_r \in \mathbb{Q}_{>0}$, and $d_1, \dots, d_r \in \mathbb{N}$ such that for any prime power q and any $b \in \mathbb{F}_q$ there exists $i \in \{1, \dots, r\}$ such that $\phi(\mathbb{F}_q; b)$ is empty or

$$|\phi(\mathbb{F}_q; b)| - \mu_i q^{d_i} < Cq^{d_i - \frac{1}{2}}. \quad (*_i)$$

Furthermore, for each $i \in \{1, \dots, r\}$ the set

$$\{b \in \mathbb{F}_q \mid (*_i) \text{ holds}\}$$

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In **red** is the monomial function (of q) that measures the approximate size of the set $\phi(\mathbb{F}_q; b)$:

$$q \mapsto \mu_i q^{d_i}.$$

Macpherson–Steinhorn:

- 1 They study classes \mathcal{C} of finite structures in which the C–D–M theorem holds (appropriately restated): *one-dimensional asymptotic classes*.
- 2 Taking an ultraproduct \mathcal{M} of such classes gives a function

$$h : \text{Def}(\mathcal{M}) \longrightarrow \mathbb{R}e^{\mathbb{N}}.$$

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They axiomatise this as follows:

Definition (Macpherson–Steinhorn, slightly reformulated)

The structure \mathcal{M} is *measurable* if there is

$$h : \text{Def}(\mathcal{M}) \longrightarrow \mathbb{R}e^{\mathbb{N}}$$

such that

- ① h is finitely additive, and if X is a singleton then $h(X) = 1e^0$;
- ② **‘MAC’ condition:** For every $\phi(x; y)$,
 - ① $\{h(\phi(\mathcal{M}; b)) \mid b \in \mathcal{M}\} = \{r_1e^{d_1}, \dots, r_ne^{d_n}\}$ is a finite set (**picture**),
and
 - ② $\{b \in \mathcal{M} \mid h(\phi(\mathcal{M}; b)) = r_ie^{d_i}\}$ is \emptyset -definable, for each i ; and
- ③ **Fubini:** If $p : X \longrightarrow Y$ is a definable surjection and h takes the constant value re^d on each fibre of p , then

$$h(X) = h(Y) \cdot re^d.$$

Theorem (Macpherson–Steinhorn)

For each pseudofinite field F we have a measuring-function

$$h_F : \text{Def}(F) \longrightarrow \mathbb{R}e^{\mathbb{N}}.$$

Proof.

Pseudofinite fields are elementarily equivalent to non-principal ultraproducts of finite fields. □

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Theorem (Macpherson–Steinhorn)

For the random graph \mathcal{R} we have a measuring-function

$$h_{\mathcal{R}} : \text{Def}(\mathcal{R}) \longrightarrow \mathbb{R}e^{\mathbb{N}}.$$

Proof.

Goes via Paley graphs. These are graphs defined on finite fields \mathbb{F}_q with $q \equiv 1 \pmod{4}$: define aEb iff $a - b$ is a square. □

Question

For the generic triangle-free graph \mathcal{H}_3 , does there exist a measuring-function

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Answer: **No**. In any measurable structure, 'dimension' is well-founded (model-theoretic language: measurable implies supersimple). However, in \mathcal{H}_3 there are infinite descending chains of definable sets with strictly decreasing dimension (model-theoretic language: dividing, tree property of the first kind).

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So it all comes to down to *dimension*.

Definition

$T = (T, +, \cdot, 0, 1, \leq)$ is a *measuring semiring* if

- 1 $(T, +, 0)$ and $(T, \cdot, 1)$ are commutative monoids,
- 2 \cdot distributes over $+$,
- 3 $(T, \leq, 0)$ is a totally ordered set with least element 0,
- 4 $\forall x, y, z (x \leq y \longrightarrow x + z \leq y + z)$,
- 5 $\forall x, y, z (x \leq y \longrightarrow x \cdot z \leq y \cdot z)$,
- 6 $\forall x 0 \cdot x = x \cdot 0 = 0$;

and

- 7 $\forall x, y, z$ if $x < y$ and (either $y \leq z \leq ny$ or $z \leq y \leq nz$, for some $n \in \mathbb{N}$) then

$$x + z < y + z.$$

Let $T = (T, +, \cdot, 0, 1, \leq)$ be a measuring semiring.

Definition (A.–Macpherson–Steinhorn–Wolf)

The structure \mathcal{M} is *T -measurable* if there is a *measuring-function*

$$h : \text{Def}(\mathcal{M}) \longrightarrow T$$

such that

- ① h is finitely additive, and if X is a singleton then $h(X) = 1$;
- ② **'MAC' condition:** For every $\phi(x; y)$,
 - ① $\{h(\phi(\mathcal{M}; b)) \mid b \in \mathcal{M}\} = \{t_1, \dots, t_n\}$ is a finite set, and
 - ② $\{b \in \mathcal{M} \mid h(\phi(\mathcal{M}; b)) = t_i\}$ is \emptyset -definable, for each i ; and
- ③ **Fubini:** If $p : X \longrightarrow Y$ is a definable surjection and h takes the constant value $t \in T$ on each fibre of p , then

$$h(X) = h(Y) \cdot t.$$

Definition (AMSW)

\mathcal{M} is *generalised measurable* if it is T -measurable for some measuring semiring T .

Free homogeneous structures

Fix a finite relational language $\mathcal{L} := \{R_1, \dots, R_r\}$.

Definition (**Free amalgamation class**)

A class \mathcal{C} of finite \mathcal{L} -structures is a *free amalgamation class* if

- 1 closed under isomorphism and substructure,
- 2 has the joint embedding property, and
- 3 has the *free amalgamation property*.

A *free homogeneous* structure is a Fraïssé amalgam of a free amalgamation class.

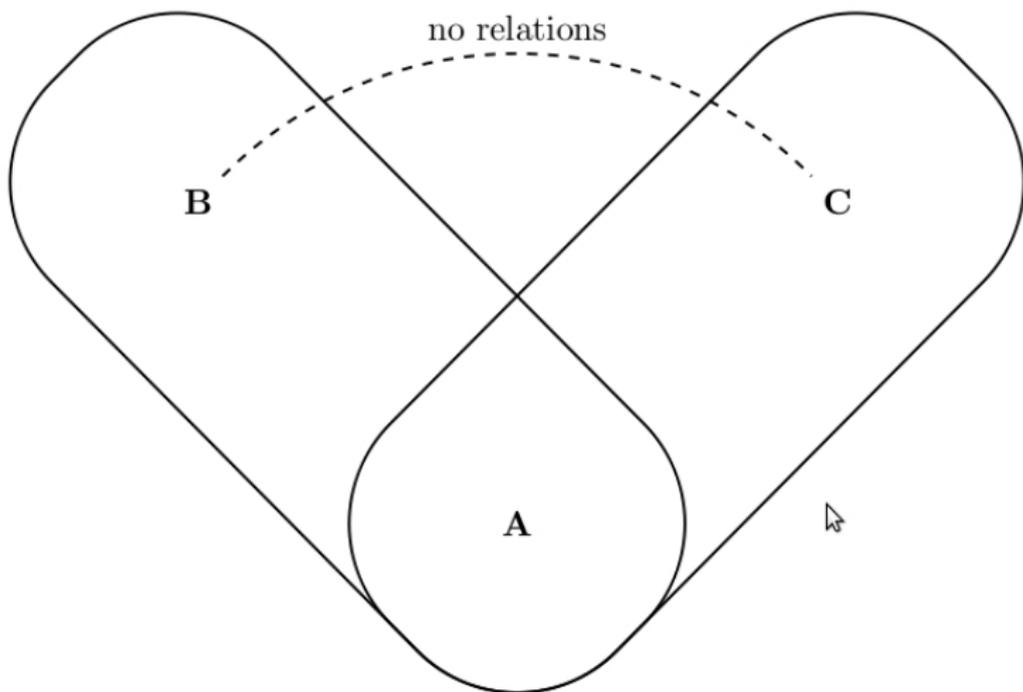


Figure: Free amalgamation

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Facts

If \mathcal{M} is the Fraïssé amalgam of the free amalgamation class \mathcal{C} then $\mathcal{C} = \text{age}(\mathcal{M})$, \mathcal{M} is \aleph_0 -categorical, (ultra)homogeneous, has quantifier elimination, and algebraic closure is trivial.

Examples

- ① $\mathcal{C}_0 := \{\text{finite graphs}\}$, $\mathcal{R} = \text{random graph}$.
- ② $\mathcal{C}_n := \{\text{finite } K_n\text{-free graphs}\}$, $\mathcal{H}_n = \text{generic } K_n\text{-free graph}$.
- ③ Hypergraph versions of these.

Non-examples

- ① class = finite partial orders, limit = generic partial order.
- ② class = finite total orders, limit = $(\mathbb{Q}, <)$.
- ③ class = finite totally ordered graphs, limit = generic totally ordered graph.

From now on we work with a fixed free homogeneous \mathcal{L} -structure \mathcal{M} .

We can think of $\mathcal{M} = \mathcal{H}_3 =$ generic triangle-free graph.

Definition of Γ_r

Definition

The structure $\Gamma_r = (\Gamma_r, +, 0, -\infty, \leq)$ is defined as follows.

- 1 Γ_r is the set

$$\Gamma_r^* := \omega^* \oplus_{\text{lex}} (-\omega) \oplus_{\text{lex}} \dots \oplus_{\text{lex}} (-\omega)$$

adjoined by two elements $-\infty$ and 0 ;

- 2 the addition $+$ is coordinate-wise on the lexicographic product, 0 is the identity, and $-\infty$ is a zero; and
- 3 \leq is such that $-\infty < 0 < \Gamma_r^*$.

Example

$$\Gamma_1 = \{-\infty\} \sqsubset_{<} \{0\} \sqsubset_{<} \left(\omega^* \oplus_{\text{lex}} (-\omega) \right).$$

Definition

Let $T_r = (T_r, \oplus, \otimes, 0e^{-\infty}, 1e^0, \leq) := \mathbb{N}e^{\Gamma_r}$.

Underlying set:

$$\{me^d \mid m \in \mathbb{N} \setminus \{0\}, d \in \Gamma_r \setminus \{-\infty\}\} \sqcup \{0e^{-\infty}\}$$

- ① $m_1e^{d_1} \oplus m_2e^{d_2} := \begin{cases} (m_1 + m_2)e^{d_1} & d_1 = d_2 \\ m_1e^{d_1} & d_1 > d_2 \\ m_2e^{d_2} & d_1 < d_2; \end{cases}$
- ② $m_1e^{d_1} \otimes m_2e^{d_2} := (m_1 \cdot m_2)e^{d_1+d_2};$
- ③ $m_1e^{d_1} \leq m_2e^{d_2}$ iff (by def.) $d_1 < d_2$ OR $(d_1 = d_2$ AND $m_1 \leq m_2)$;
- ④ $0e^{-\infty}$ is the \oplus -identity and a ' \otimes -zero', and $1e^0$ is the \otimes -identity.

Fact

T_r is a measuring semiring.

Theorem (Main Theorem again)

\mathcal{M} is T_r -measurable with measuring-function $h_{\mathcal{M}}$.

Finishing the proof

It remains to argue that

$$h_{\mathcal{M}} : \text{Def}(\mathcal{M}) \longrightarrow T_r$$

really is a measuring-function.

- 1 Finitely additive
- 2 Finite sets: a single tuple is a complete type
- 3 'mac' condition: follows from \aleph_0 -categoricity.
- 4 **Fubini condition:**