

Mixing time for a random walk on a ring

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Joint work with Michael Bate

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Aspects of Random Walks

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Random walks on groups

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For example, we could have $G = S_n$, the symmetric group on $\{1, 2, \dots, n\}$, and we could set

$$P(g) = \begin{cases} \frac{1}{n} & \text{if } g = 1 \text{ is the identity} \\ \frac{2}{n^2} & \text{if } g = (i, j) \text{ is a transposition} \\ 0 & \text{otherwise} \end{cases}$$

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A *random walk on G* is then a Markov chain X with transitions governed by the distribution P . So we fix a starting point X_0 , and then set

$$\mathbb{P}(X_{t+1} = hg \mid X_t = g) = P(h)$$

Distribution after repeated steps is given by *convolution*:

$$\mathbb{P}(X_2 = g | X_0 = 1) = P * P(g) = \sum_h P(gh^{-1})P(h)$$

As long as the probability distribution P isn't concentrated on a subgroup, the stationary distribution π for X is the uniform distribution; $\pi(g) = 1/|G|$ for all $g \in G$. When X is ergodic, we're interested in how long it takes for it to converge to equilibrium.

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Definition

The **mixing time** of X is

$$\tau(\varepsilon) = \min \{t : \|\mathbb{P}(X_t \in \cdot) - \pi(\cdot)\|_{TV} \leq \varepsilon\} .$$

Cutoff phenomenon

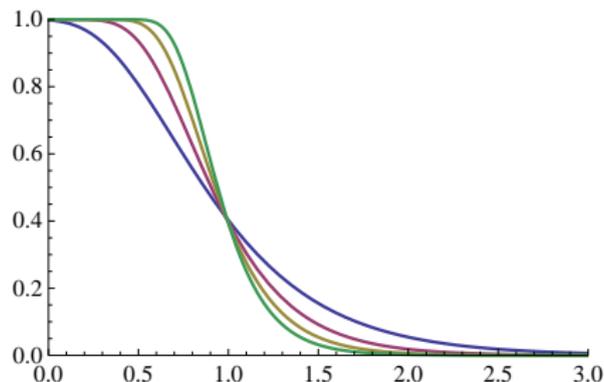
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In lots of nice examples a **cutoff phenomenon** is exhibited:

$$\text{for all } \varepsilon > 0, \quad \lim_{n \rightarrow \infty} \frac{\tau^{(n)}(\varepsilon)}{\tau^{(n)}(1 - \varepsilon)} = 1$$



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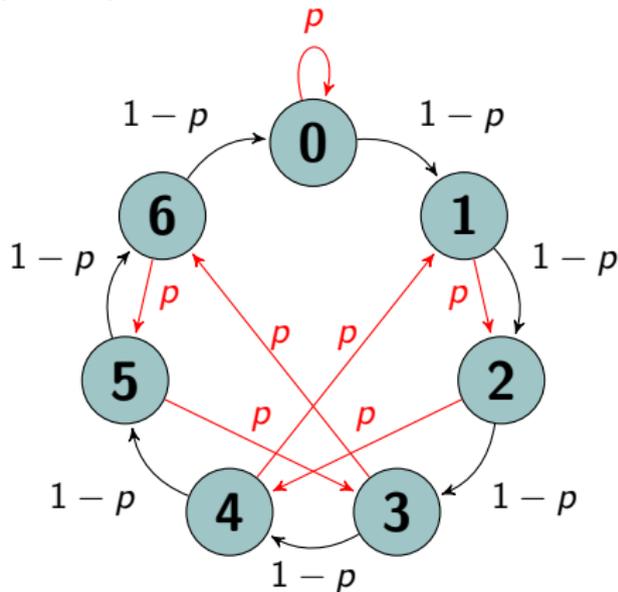
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Related results

Chung, Diaconis and Graham (1987) study the process (used in random number generation)

$$x \rightarrow \begin{cases} ax - 1 & \text{w.p. } \frac{1}{3} \\ ax & \text{w.p. } \frac{1}{3} \\ ax + 1 & \text{w.p. } \frac{1}{3} \end{cases}$$

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- When $a = 2$ and $n = 2^m - 1$ there exist constants c and c' such that:

$$\text{for } t_n \geq c \log n \log \log n, \quad \left\| P_{t_n}^{(n)} - \pi^{(n)} \right\|_{\text{TV}} \rightarrow 0 \text{ as } n \rightarrow \infty$$

$$\text{for } t_n \leq c' \log n \log \log n, \quad \left\| P_{t_n}^{(n)} - \pi^{(n)} \right\|_{\text{TV}} > \varepsilon \text{ as } n \rightarrow \infty$$

Back to our process...

General problem

The distribution of X_t isn't given by convolution.

$$X_t = 2I_t X_{t-1} + (1 - I_t)(X_{t-1} + 1)$$

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So consider (with $X_0 = Y_0 = 0$)

$$Y_k = \sum_{j=1}^k 2^{k+1-j} S_j \pmod{n}, \quad S_j \stackrel{\text{i.i.d.}}{\sim} \text{Geom}(p_n) \pmod{n}$$

(i.e. Y_k is the position of X immediately following the k^{th} jump.)

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Restrict attention to $p_n = \frac{1}{2n^\alpha}$, $\alpha \in (0, 1)$.

We expect things to happen (for Y) sometime around

$$T_n := \log_2(np_n) \sim (1 - \alpha) \log_2 n$$

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Theorem

Y exhibits a **cutoff** at time T_n , with window size $O(1)$.

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To prove this, we need to show that (for n odd)

$$\liminf_{n \rightarrow \infty} \left\| \mathbb{P} \left(X_{T_n - \theta}^{(n)} \in \cdot \right) - \frac{1}{n} \right\|_{\text{TV}} \geq 1 - \varepsilon(\theta)$$

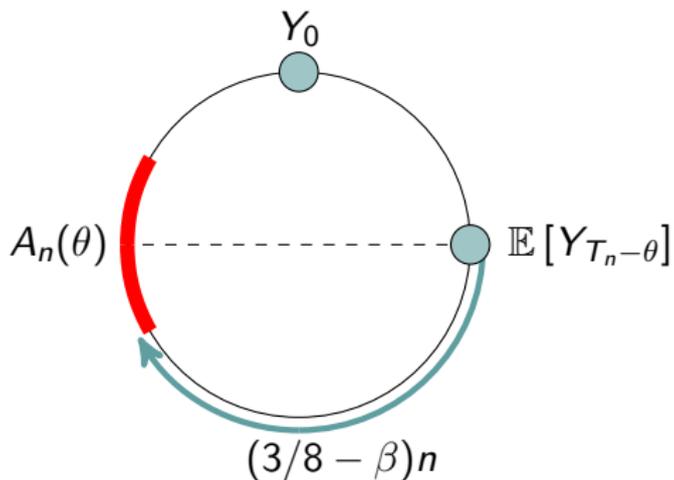
and

$$\limsup_{n \rightarrow \infty} \left\| \mathbb{P} \left(X_{T_n + \theta}^{(n)} \in \cdot \right) - \frac{1}{n} \right\|_{\text{TV}} \leq \varepsilon(\theta),$$

where $\varepsilon(\theta) \rightarrow 0$ as $\theta \rightarrow \infty$.

Lower bound

For a lower bound, we simply find a set $A_n(\theta)$ (of considerable size) that Y has very little chance of hitting before time $T_n - \theta$, for large $\theta \in \mathbb{N}$. (Recall the definition of total variation distance!)



- If $A_n(\theta)$ is chosen as above then $\pi(A_n(\theta)) = 1/4 + 2\beta$.

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Lemma (Lower bound for Y)

For $\theta \geq 3$,

$$\|\mathbb{P}(Y_{T_n-\theta} \in \cdot) - \pi_n(\cdot)\|_{\text{TV}} \geq 1 - 4^{1-\theta/3}.$$

Upper bound

Let P be a probability on a group G . A (complex) representation ρ is a group homomorphism $\rho : G \rightarrow GL_d(\mathbb{C})$, where $GL_d(\mathbb{C})$ denotes the group of $d \times d$ invertible complex matrices. We write

$$\hat{P}(\rho) = \sum_{g \in G} P(g)\rho(g)$$

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This behaves very nicely with respect to convolution:

$$\widehat{P * P}(\rho) = \hat{P}(\rho)\hat{P}(\rho)$$



A basic but extremely useful result is the following:

Lemma (Diaconis and Shahshahani, 1981)

$$\|P - \pi\|_{\text{TV}}^2 \leq \frac{1}{4} \sum_{\substack{\text{non-triv} \\ \text{irr } \rho}} d_\rho \text{tr} \left(\hat{P}(\rho) \hat{P}(\rho)^* \right)$$

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Our subsampled walk Y is a random walk on the group $(\mathbb{Z}_n, +)$, whose n irreducible (one-dimensional) representations are determined by

$$\rho_s(1) := e^{i \frac{2\pi s}{n}} \quad \text{for } 0 \leq s \leq n-1$$



The Upper Bound Lemma becomes

$$\|P - \pi\|_{\text{TV}}^2 \leq \frac{1}{4} \sum_{s=1}^{n-1} |\hat{P}(\rho_s)|^2$$

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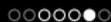
Substituting the correct distribution for Y_t leads us to the following upper bound:

$$\|\delta_0 P_t - \pi\|_{\text{TV}}^2 \leq \frac{1}{4} \sum_{s=1}^{n-1} \prod_{k=1}^t \frac{p_n^2}{1 - 2(1 - p_n) \cos(\frac{2\pi}{n} 2^k s) + (1 - p_n)^2}$$

Lemma (Upper bound for Y)

Let $p_n = 1/2n^\alpha$, with $\alpha \in (0, 1]$. For $\theta \in \mathbb{N}$,

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Proof.

Careful analysis of the right-hand side!

(Identify which terms really contribute to the sum ($s = (n \pm 1)/2$ accounts for nearly everything), deal with these, and show that nothing else really matters.) □

This completes our proof of a cutoff for Y .

Moving from Y to X

We've seen that Y mixes in an interval of length $O(1)$ around $T_n = \log_2(np_n)$: what does this tell us about the mixing time for X ?

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Corollary

For $p_n = 1/2n^\alpha$, with $\alpha \in (0, 1)$, X exhibits a cutoff at time

$$T_n^X = T_n/p_n = 2(1 - \alpha)n^\alpha \log_2 n$$

with window size $\sqrt{T_n}/p_n$.

Proof.

Essentially follows from the observation that the number of jumps by time $T_n^X + c\sqrt{T_n}/p_n$ concentrates (in an interval of order $\sqrt{T_n}$) around $T_n + c\sqrt{T_n}$. □

And finally: open problems

- 1 We can deal with more interesting steps in our walk, but not yet with more interesting jumps, e.g. consider

$$x \rightarrow \begin{cases} x + 1 & \text{w.p. } 1 - p_n \\ 2x & \text{w.p. } p_n/2 \\ \left(\frac{n+1}{2}\right)x & \text{w.p. } p_n/2 \end{cases}$$

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- ② Or how about this process?

$$x \rightarrow \begin{cases} x + 1 & \text{w.p. } 1 - p_n \\ ax & \text{w.p. } p_n \end{cases}$$

where multiplication by a is not invertible? (Stationary distribution won't even be uniform...)