

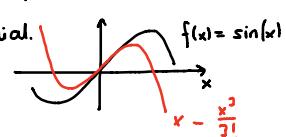
APPROXIMATION THEORY

key idea: Replace a complicated function by a simple function.

e.g. 1). $\pi \rightarrow 3.14$.

2). function $f(x) \rightarrow$ tangent $f'(x)$

3). function $f(x) \rightarrow$ Taylor polynomial.
$$\sum_{n=0}^{\infty} \frac{f^{(n)}(0)}{n!} x^n$$



4). All of science/engineering measurements \rightarrow finite precision.

Why are simple functions better?

- Easier to store (discrete representation - e.g. mp3 or jpeg compression).
- Easier to evaluate.
- Easier to compute with (e.g. differentiate, find roots, etc.).

We aim to represent a complicated function $f(x)$ by a simpler function $\phi(x; a_0, \dots, a_n)$ where a_0, \dots, a_n are parameters chosen to give the best approximation of f .

Common ways to do this are:

1). Interpolation - choose the a_i so that

$$\phi(x_i; a_0, \dots, a_n) = f(x_i) \text{ on a prescribed set of points } \{x_i \mid i=0, \dots, n\}. \quad \leftarrow \begin{matrix} \text{seen in 2H Num. An.} \\ (\text{will study more}) \end{matrix}$$

2). Least-squares approximation - choose the a_i to minimize

$$\|f - \phi\|_2 := \left(\int_a^b (f(x) - \phi(x; a_0, \dots, a_n))^2 dx \right)^{1/2}. \quad \leftarrow \begin{matrix} \text{seen in 2H Num. An. (not in this course)} \end{matrix}$$

3). Minimax approximation - choose the a_i to minimize

$$\|f - \phi\|_\infty := \max_{x \in [a, b]} |f(x) - \phi(x; a_0, \dots, a_n)|. \quad \leftarrow \begin{matrix} \text{will study.} \end{matrix}$$

COURSE OUTLINE

- 0). Revision of polynomial interpolation.
- 1). Piecewise polynomial interpolation (splines).
- 2). Minimax approximation.
- 3). Trigonometric interpolation. (discrete Fourier transform).

0) REVISION OF POLYNOMIAL INTERPOLATION. (2H Numerical Analysis).

- Main points:-
- (1) Given $f \in C[a,b]$ and $n+1$ distinct points $a \leq x_0 < x_1 < \dots < x_n \leq b$, there is a unique polynomial $p_n \in P_n$ such that $p_n(x_i) = f(x_i)$ for $i=0, \dots, n$.
polynomial of degree $\leq n$.
 - (2) You have to be very careful where the points x_i are placed.

The proof of existence is by construction with Lagrange polynomials.

e.g. $n=1$

$$p_1(x) = \frac{x - x_0}{x_1 - x_0} f_0 + \frac{x - x_1}{x_1 - x_0} f_1.$$

In general,

$$p_n(x) = \sum_{i=0}^n f_i l_i(x) \quad \text{where} \quad l_i(x) = \prod_{\substack{j=0 \\ j \neq i}}^{n-1} \frac{x - x_j}{x_i - x_j}$$

are the Lagrange polynomials.

You should have seen the error formula:

Thm - Let $f \in C^{n+1}[a,b]$ and $a \leq x_0 < x_1 < \dots < x_n \leq b$. Let $p_n \in P_n$ be such that $p_n(x_i) = f(x_i)$ for $i=0, \dots, n$.

O.1 Then for every $x \in [a,b]$ there is a $\xi \in [a,b]$, depending on x , such that

$$f(x) - p_n(x) = \frac{W_{n+1}(x) f^{(n+1)}(\xi)}{(n+1)!}$$

where $W_{n+1}(x) = \prod_{j=0}^n (x - x_j)$ is the error polynomial.

- Disadvantage: can only apply this formula if f is smooth enough.

- Can use this to put an upper bound:

$$\|f - p_n\|_\infty \leq \frac{\|W_{n+1}\|_\infty \|f^{(n+1)}\|_\infty}{(n+1)!}$$

The ∞ -norm means $\|f\|_\infty = \max_{x \in [a,b]} |f(x)|$.

Suppose we take equally-spaced points $x_j = a + \left(\frac{b-a}{n}\right)j$ for $j=0, \dots, n$. As we increase the resolution n , we would like $\|f - p_n\|_\infty \rightarrow 0$ as $n \rightarrow \infty$.

Problem- The sequence $\{p_n\}$ may not converge to f because $\|W_{n+1}\|_\infty \|f^{(n+1)}\|_\infty$ may tend to ∞ faster than $(n+1)!$.

A famous example is due to Runge (1901):

Example: $f(x) = \frac{1}{1+25x^2} \quad x \in [-1, 1]$.

This function is well-behaved with all derivatives continuous and bounded, yet p_n diverges for equally-spaced points.

Remark- Actually it does converge near $x=0$. These kind of interpolants are used in higher-order (e.g. 6th) finite difference methods for solving PDEs.

